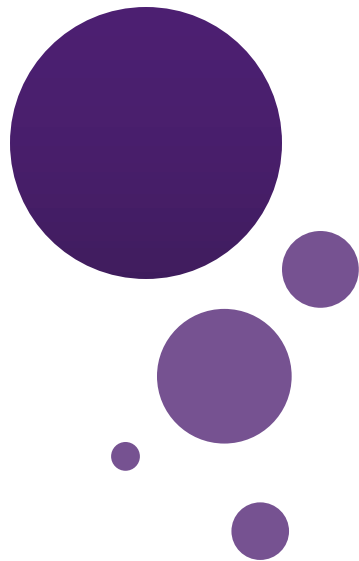




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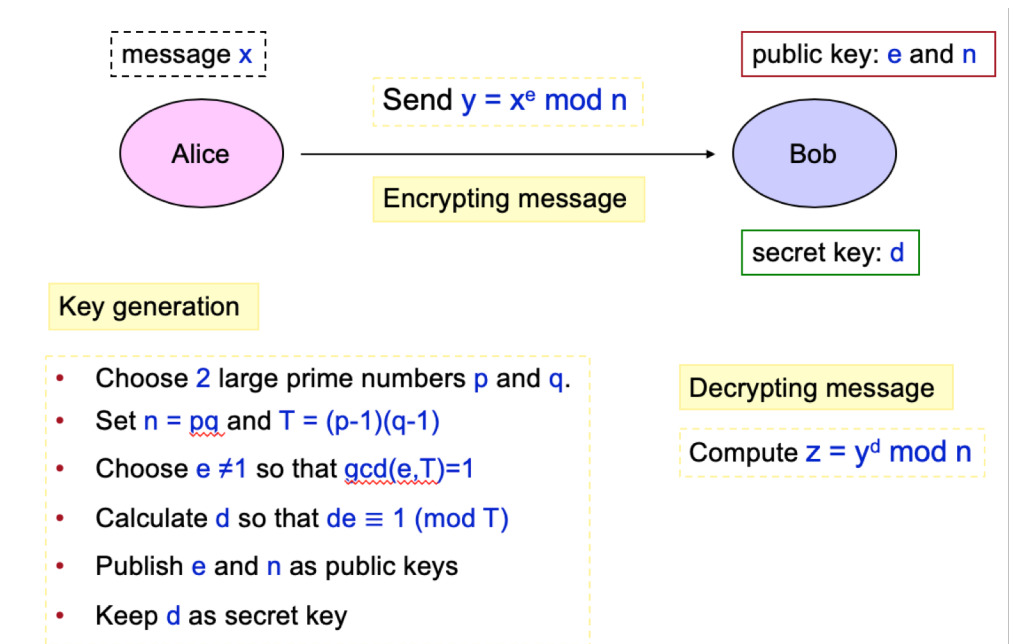
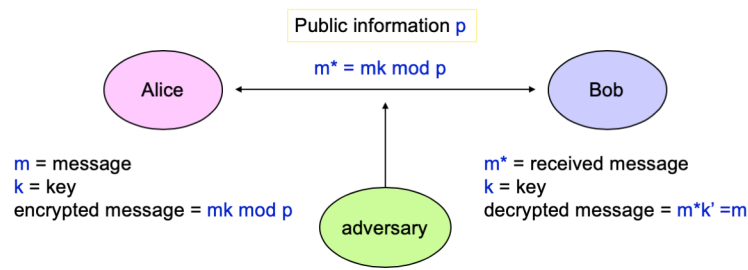
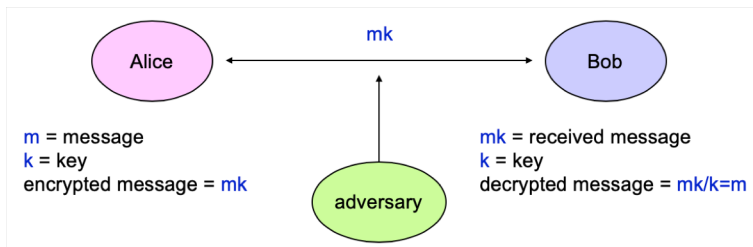


Lecture 21: Mathematic Induction

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Recap Previous Lecture

- Turing Code Ver 1 & 2.
- RSA Cryptosystem



Outline

- The Ideal of Mathematic Induction
- Basic Induction Proofs
- Applications
- A Paradox

Outline

- **The Ideal of Mathematic Induction**
- Basic Induction Proofs
- Applications
- A Paradox

Odd Power are Odd

Fact: If m is odd and n is odd, then nm is odd.

Proposition: for an odd number m , m^k is odd for all non-negative integer k .

$$\forall k \in \mathbb{N} \text{ odd}(m^k)$$

Let $P(i)$ be the proposition that m^i is odd.

$$\forall k \in \mathbb{N} P(k)$$

Idea of induction.

- $P(1)$ is true by definition.
- $P(2)$ is true by $P(1)$ and the fact.
- $P(3)$ is true by $P(2)$ and the fact.
- $P(i+1)$ is true by $P(i)$ and the fact.
- So $P(i)$ is true for all i .

Divisibility by a Prime

Theorem. Any integer $n > 1$ is divisible by a prime number.

- Let n be an integer.
- If n is a prime number, then we are done.
- Otherwise, $n = ab$, both are smaller than n .
- If a or b is a prime number, then we are done.
- Otherwise, $a = cd$, both are smaller than a .
- If c or d is a prime number, then we are done.
- Otherwise, repeat this argument, since the numbers are getting smaller and smaller, this will eventually stop and we have found a prime factor of n .

Idea of induction.

Idea of Induction

Objective: Prove $\forall n \geq 0 P(n)$

This is to prove

$$\underline{P(0)} \wedge \underline{P(1)} \wedge \underline{P(2)} \wedge \dots \wedge \underline{P(n)} \dots$$

The idea of induction is to first prove $P(0)$ unconditionally,
then use $P(0)$ to prove $P(1)$
then use $P(1)$ to prove $P(2)$
and repeat this to infinity...

Idea of Induction

0 and (from n to $n+1$),

proves 0, 1, 2, 3,....

Very easy
to prove

Much easier to
prove with $P(n)$
as an assumption.

$$P(0), P(n) \rightarrow P(n+1)$$

$$\forall m \in \mathbb{N}. P(m)$$

For any $n \geq 0$

Like domino effect...



Outline

- The Ideal of Mathematic Induction
- **Basic Induction Proofs**
- Applications
- A Paradox

Proof by Induction

Let's prove:

$$\forall r \neq 1. 1 + r + r^2 + \dots + r^n = \frac{r^{n+1} - 1}{r - 1}$$

Statements in **green** form a template for inductive proofs.

Proof: (by induction on n)

The induction hypothesis, $P(n)$, is:

$$\forall r \neq 1. 1 + r + r^2 + \dots + r^n = \frac{r^{n+1} - 1}{r - 1}$$

Proof by Induction

Induction Step: Assume $P(n)$ for some $n \geq 0$ and prove $P(n + 1)$:

$$\forall r \neq 1. 1 + r + r^2 + \cdots + r^{n+1} = \frac{r^{(n+1)+1} - 1}{r - 1}$$

Have $P(n)$ by assumption:

So let r be any number $\neq 1$, then from $P(n)$ we have

$$1 + r + r^2 + \cdots + r^n = \frac{r^{n+1} - 1}{r - 1}$$

How do we proceed?

Proof by Induction

adding r^{n+1} to both sides,

$$\begin{aligned}1 + \cdots + r^n + r^{n+1} &= \frac{r^{n+1} - 1}{r - 1} + r^{n+1} \\ &= \frac{r^{n+1} - 1 + r^{n+1}(r - 1)}{r - 1} \\ &= \frac{r^{(n+1)+1} - 1}{r - 1}\end{aligned}$$

$$\forall r \neq 1. \quad 1 + r + r^2 + \cdots + r^{n+1} = \frac{r^{(n+1)+1} - 1}{r - 1}$$

which is $P(n+1)$. This completes the induction proof.

Proving an Equality

$$\forall n \geq 1 \quad 1^3 + 2^3 + \dots + n^3 = \left(\frac{n(n+1)}{2}\right)^2$$

Let $P(n)$ be the induction hypothesis that the statement is true for n .

Base case: $P(1)$ is true

Induction step: assume $P(n)$ is true, prove $P(n+1)$ is true.

$$\begin{aligned} & 1^3 + 2^3 + \dots + n^3 + (n+1)^3 \\ &= \left(\frac{n(n+1)}{2}\right)^2 + (n+1)^3 \quad \text{by induction} \\ &= (n+1)^2(n^2/4 + n + 1) \\ &= (n+1)^2\left(\frac{n^2 + 4n + 4}{4}\right) = \left(\frac{(n+1)(n+2)}{2}\right)^2 \end{aligned}$$

Proving a Property

$\forall n \geq 1, 2^{2n} - 1$ is divisible by 3

Base Case ($n = 1$): $2^{2n} - 1 = 2^2 - 1 = 3$

Induction Step: Assume $P(i)$ for some $i \geq 1$ and prove $P(i + 1)$:

Assume $2^{2i} - 1$ is divisible by 3, prove $2^{2(i+1)} - 1$ is divisible by 3.

$$\begin{aligned} 2^{2(i+1)} - 1 &= 2^{2i+2} - 1 \\ &= 4 \cdot 2^{2i} - 1 \\ &= \underbrace{3 \cdot 2^{2i}} + \underbrace{2^{2i} - 1} \end{aligned}$$

Divisible by 3

Divisible by 3 by induction

Proving a Property

$$\forall n \geq 2, \quad n^3 - n \text{ is divisible by } 6$$

Base Case ($n = 2$): $2^3 - 2 = 6$

Induction Step: Assume $P(i)$ for some $i \geq 2$ and prove $P(i + 1)$:

Assume $n^3 - n$ is divisible by 6

Prove $(n + 1)^3 - (n + 1)$ is divisible by 6.

$$\begin{aligned}(n + 1)^3 - (n + 1) &= (n^3 + 3n^2 + 3n + 1) - (n + 1) \\ &= \underbrace{(n^3 - n)}_{\text{Divisible by 6 by induction}} + 3 \underbrace{(n^2 + n)}_{\text{Divisible by 2 by case analysis}}\end{aligned}$$

Divisible by 6
by induction

Divisible by 2
by case analysis

Proving an Inequality

$$\forall n \geq 3, \quad 2n + 1 < 2^n$$

Base Case ($n = 3$): $2n + 1 = 7 < 2^n = 2^3 = 8$

Induction Step: Assume $P(i)$ for some $i \geq 3$ and prove $P(i + 1)$:

Assume $2i + 1 < 2^i$, prove $2(i + 1) + 1 < 2^{(i+1)}$

$$2(i + 1) + 1 = 2i + 1 + 2$$

$$< 2^i + 2 \quad \text{by induction}$$

$$< 2^i + 2^i \quad \text{since } i \geq 3$$

$$= 2^{(i+1)}$$

Proving an Inequality

$$\forall n \geq 2, \quad \frac{1}{\sqrt{1}} + \frac{1}{\sqrt{2}} + \dots + \frac{1}{\sqrt{n}} > \sqrt{n}$$

Base Case ($n = 2$): is true

Induction Step: Assume $P(i)$ for some $i \geq 2$ and prove $P(i + 1)$:

$$\begin{aligned} & \frac{1}{\sqrt{1}} + \frac{1}{\sqrt{2}} + \dots + \frac{1}{\sqrt{n}} + \frac{1}{\sqrt{n+1}} \\ > \sqrt{n} + \frac{1}{\sqrt{n+1}} & \text{by induction} \\ = \frac{\sqrt{n}\sqrt{n+1} + 1}{\sqrt{n+1}} \\ > \frac{\sqrt{n}\sqrt{n} + 1}{\sqrt{n+1}} = \frac{n+1}{\sqrt{n+1}} \\ = \sqrt{n+1} \end{aligned}$$

Outline

- The Ideal of Mathematic Induction
- Basic Induction Proofs
- **Applications**
- A Paradox

Gray Code

Can you find an ordering of all the n -bit strings in such a way that two consecutive n -bit strings differed by only one bit?

This is called the Gray code and has many applications.

How to construct them?

Think inductively! (or recursively!)

2 bit

3 bit

00

000

01

001

11

011

10

010

110

111

101

100

Can you see the pattern?

How to construct 4-bit gray code?

Gray Code

3 bit	3 bit (reversed)
000	100
001	101
011	111
010	110
110	010
111	011
101	001
100	000

Every 4-bit string appears exactly once.

4 bit

0000
0001
0011 ← differed by 1 bit
0010 ← by induction
0110
0111
0101
0100 ← differed by 1 bit
1100 ← by construction
1101
1111
1110
1010 ← differed by 1 bit
1011 ← by induction
1001
1000

Gray Code

n bit	n bit (reversed)
000...0	100...0
...	...
...	...
...	...
...	...
...	...
...	...
...	...
100...0	000...0

n+1 bit

0 000...0

0 ...

0 ...

0 ...

0 ...

0 ...

0 ...

0 100...0

1 100...0

1 ...

1 ...

1 ...

1 ...

1 ...

1 ...

1 000...0

differed by 1 bit
by induction

differed by 1 bit
by construction

differed by 1 bit
by induction

Every (n+1)-bit string appears exactly once.

So, by induction,
Gray code exists for any n.

Hadamard Matrix (Optional)

Can you construct an $n \times n$ matrix with all entries ± 1 and all the rows are orthogonal to each other?

Two rows are orthogonal if their inner product is zero.

That is, let $a = (a_1, \dots, a_n)$ and $b = (b_1, \dots, b_n)$,

their inner product $ab = a_1b_1 + a_2b_2 + \dots + a_nb_n$

This matrix is famous and has many applications.

To think inductively, first we come up with small examples.

$$\begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

Hadamard Matrix (Optional)

Then we use the small examples to build larger examples.

Suppose we have an $n \times n$ Hadamard matrix H_n .

We can use it to construct an $2n \times 2n$ Hadamard matrix as follows.

$$\begin{bmatrix} H_n & H_n \\ H_n & -H_n \end{bmatrix}$$

Check this!

So by induction there is a $2^k \times 2^k$ Hadamard matrix for any k .

Inductive Construction

This technique is very useful.

We can use it to construct:

- codes
- graphs
- matrices
- circuits
- algorithms
- designs
- proofs
- buildings
- ...

Outline

- The Ideal of Mathematic Induction
- Basic Induction Proofs
- An Interesting Example
- **A Paradox**

Paradox

Theorem: All horses are the same color.

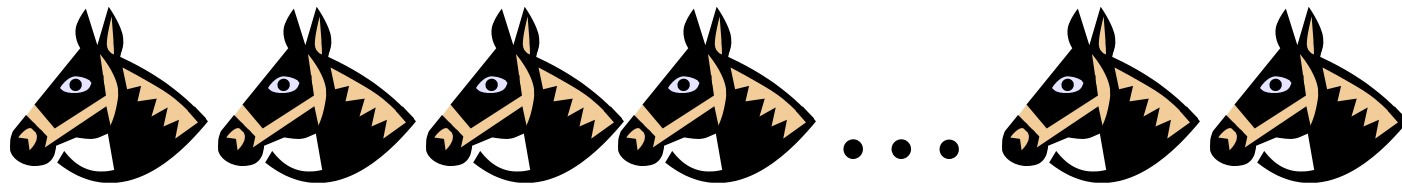
Proof: (by induction on n)

Induction hypothesis:

$P(n) ::=$ any set of n horses have the same color

Base case ($n=0$):

No horses so *obviously* true!

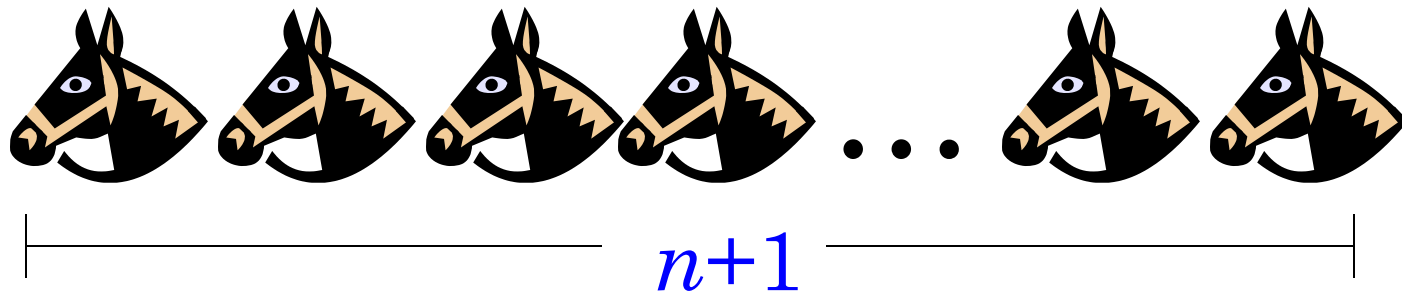


Paradox

(Inductive case)

Assume any n horses have the same color.

Prove that any $n+1$ horses have the same color.

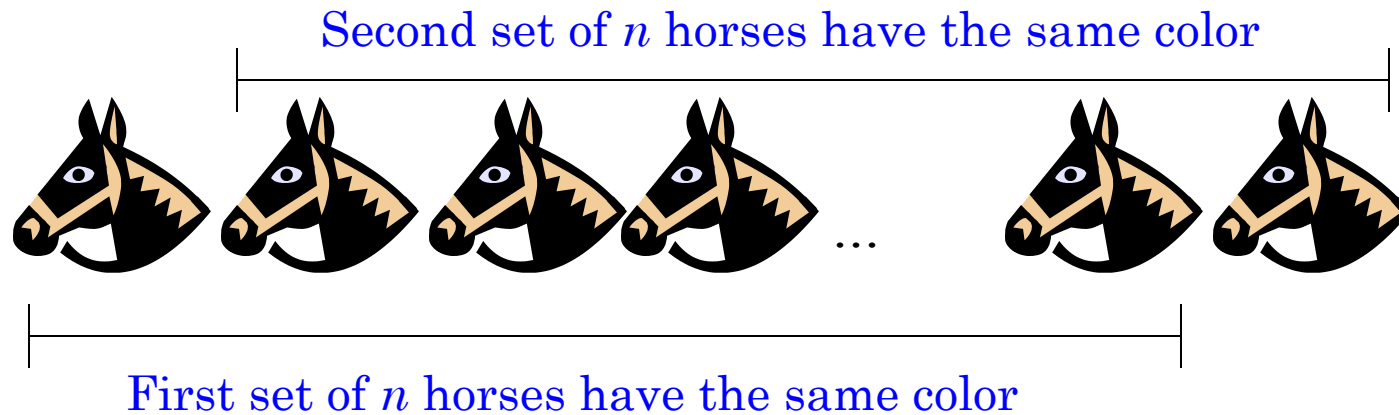


Paradox

(Inductive case)

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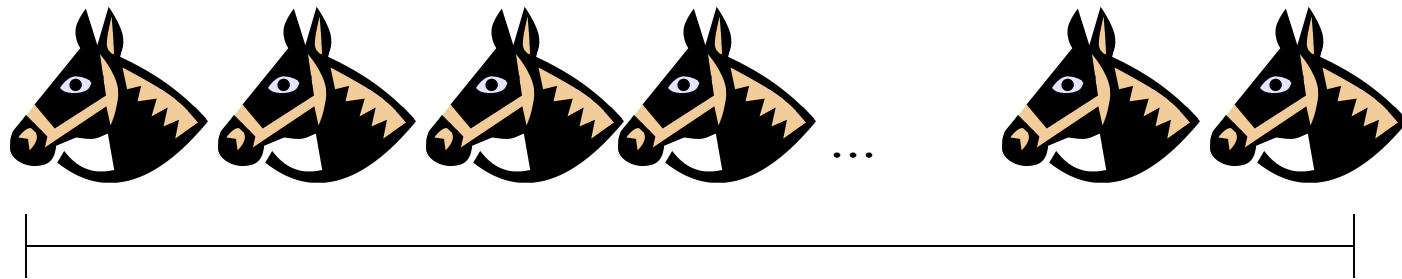


Paradox

(Inductive case)

Assume any n horses have the same color.

Prove that any $n+1$ horses have the same color.

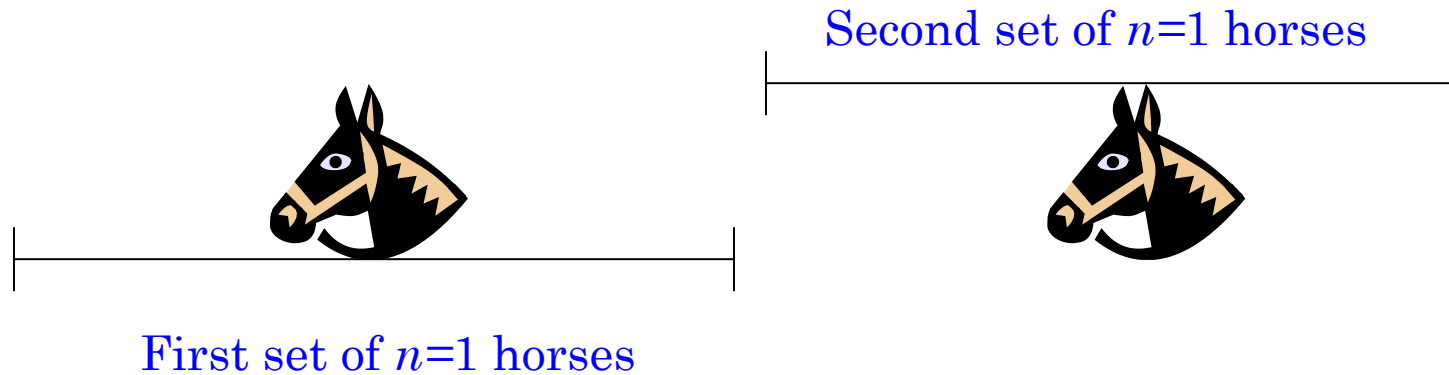


Therefore the set of $n+1$ have the same color!

Paradox

What is wrong? $n = 1$

Proof that $P(n) \rightarrow P(n+1)$
is false if $n = 1$, because the two
horse groups *do not overlap*.



(But proof works for all $n \neq 1$)

Summary

You should understand the principle of mathematical induction well, and do basic induction proofs like

- proving equality
- proving inequality
- proving property

Mathematical induction has a wide range of applications in computer science.

In the next lecture we will see more applications and more techniques.

Next class

- Topic: Strong Induction and Well-Ordering
- Pre-class reading: Chap 5.2

