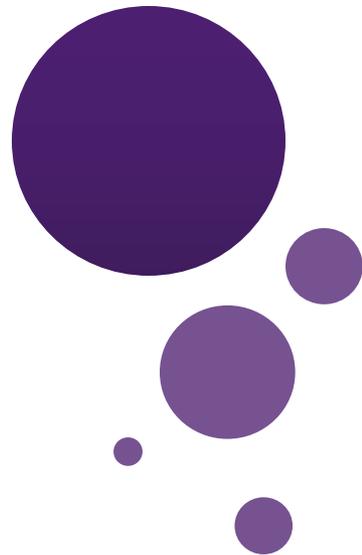




UNIVERSITY
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Lecture 19: Primes and Greatest Common Divisors



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About the Midterm Exam 1

University at Albany, SUNY

College of Engineering and Applied Sciences, Computer Science

ISEN/ISCI-210: Discrete Structures

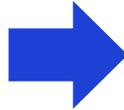
Fall 2018

Midterm Exam 1

Name: _____ ID #: _____ Score: _____

- This is a CLOSE BOOK & CLOSE NOTE exam. Also, you cannot access the Internet or use your laptop computer. Do the exam independently.
- Logical equivalence tables are given on Page 2.
- There are a total of 100 points in the exam. Plan your work accordingly.
- Write out the steps for all problems to receive the full credit. Use additional pages if necessary.
- Date: Oct 8th, 2018.
- Location: Lecture center hall 25.
- Time: 9:20 am - 10:20 am (can be extended to 10:35 am).

Problem	Points	Scores
Problem 1: True or False	20	
Problem 2: True Table and Logical Equivalence	20	
Problem 3: Predicatives and Quantifiers	20	
Problem 4: Set, Sequences and Summation	20	
Problem 5: Functions	20	



University at Albany, SUNY

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Spring 2019

Midterm Exam 1

Name: _____ ID #: _____ Score: _____

- This is a CLOSE BOOK & CLOSE NOTE exam. Also, you cannot access the Internet or use your laptop computer. Do the exam independently.
- There are a total of 100 points in the exam. Plan your work accordingly.
- Write out the steps for all problems to receive the full credit. Use additional pages if necessary.
- Date: March 8th, 2019.
- Location: Lecture center hall 25.
- Time: 11:30 am - 12:25 pm.

Problem	Points	Scores
Problem 1: Logical Equivalence	20	
Problem 2: Sequence and Summation	30	
Problem 3: Functions	30	
Problem 4: Algorithm, Growth Function and Complexity	20	

Attentions: No make-up tests or remote tests are offered this semester.

Recap Previous Lecture

- Integer Representation, Base Conversion.
- Application: Binary Modular Exponentiation.

$175^{235} \bmod 257$

- $d = 175 \bmod 257$
- $d = 175^2 \times 175 \bmod 257$
- $d = 175^8 \times 175^2 \times 175 \bmod 257$
- $d = 175^{32} \times 175^8 \times 175^2 \times 175 \bmod 257$
- $d = 175^{64} \times 175^{32} \times 175^8 \times 175^2 \times 175 \bmod 257$
- $d = 175^{128} \times 175^{64} \times 175^{32} \times 175^8 \times 175^2 \times 175 \bmod 257$

$$d = t \times d \bmod 257 \text{ when } t=175^k \text{ and } k=2^{j-1}$$

- $t = 175^2 \bmod 257 = 42$
- $t = 175^4 \bmod 257 = (175^2 \bmod 257)^2 \bmod 257 = 222$
- $t = 175^8 \bmod 257 = (175^4 \bmod 257)^2 \bmod 257 = 197$
- $t = 175^{16} \bmod 257 = (175^8 \bmod 257)^2 \bmod 257 = 2$
- $t = 175^{32} \bmod 257 = (175^{16} \bmod 257)^2 \bmod 257 = 4$
- $t = 175^{64} \bmod 257 = (175^{32} \bmod 257)^2 \bmod 257 = 16$
- $t = 175^{128} \bmod 257 = (175^{64} \bmod 257)^2 \bmod 257 = 256$

$$t = t^2 \bmod 257$$

Outline

- Prime and Composite
- Prime Factorizations
- Distribution of Primes
- GCD and LCM
- Euclidean Algorithm

Outline

- **Prime and Composite**
- Prime Factorizations
- Distribution of Primes
- Greatest Common Divisor (GCD)
- Least Common Multiple (LCM)
- Euclidean Algorithm

Prime, Composite and Theorem 1

- **Prime:** a positive integer p greater than 1 if the only positive factors of p are 1 and p
- A positive integer greater than 1 that is not prime is called **composite**

THE FUNDAMENTAL THEOREM OF ARITHMETIC Every integer greater than 1 can be written uniquely as a prime or as the product of two or more primes where the prime factors are written in order of nondecreasing size.

Example

- Prime factorizations of integers
 - $100=2\cdot 2\cdot 5\cdot 5=2^2\cdot 5^2$
 - $641=641$
 - $999=3\cdot 3\cdot 3\cdot 37=3^3\cdot 37$
 - $1024=2\cdot 2\cdot 2\cdot 2\cdot 2\cdot 2\cdot 2\cdot 2\cdot 2\cdot 2=2^{10}$

Theorem 2

If n is a composite integer, then n has a prime divisor less than or equal to \sqrt{n} .

- As n is composite, n has a factor $1 < a < n$, and thus $n = ab$
- We show that $a \leq \sqrt{n}$ or $b \leq \sqrt{n}$ (by contraposition)
- Thus n has a divisor not exceeding \sqrt{n}
- This divisor is either prime or by the fundamental theorem of arithmetic, has a prime divisor less than itself, and thus a prime divisor less than \sqrt{n}
- In either case, n has a prime divisor $b \leq \sqrt{n}$

Example

- Show that 101 is prime
- The only primes not exceeding $\sqrt{101}$ are 2, 3, 5, 7.
- As 101 is not divisible by 2, 3, 5, 7, it follows that 101 is prime.

Outline

- Prime and Composite
- **Prime Factorizations**
- Distribution of Primes
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- Euclidean Algorithm

Procedure for prime factorization

- Begin by dividing n by successive primes, starting with 2
- If n has a prime factor, we would find a prime factor not exceeding \sqrt{n} .
- If no prime factor is found, then n is prime
- Otherwise, if a prime factor p is found, continue by factoring n/p

Procedure for prime factorization

- Note that n/p has no prime factors less than p
- If n/p has no prime factor greater than or equal to p and not exceeding its square root, then it is prime
- Otherwise, if it has a prime factor q , continue by factoring $n/(pq)$
- Continue until factorization has been reduced to a prime

Example

- Find the prime factorization of 7007
- Start with 2, 3, 5, and then 7, $7007/7=1001$
- Then, divide 1001 by successive primes, beginning with 7, and find $1001/7=143$
- Continue by dividing 143 by successive primes, starting with 7, and find $143/11=13$
- As 13 is prime, the procedure stops
- $7007=7 \cdot 7 \cdot 11 \cdot 13=7^2 \cdot 11 \cdot 13$

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Theorem 3

There are infinitely many primes.

- Proof by contradiction
- Assume that there are only finitely many primes, p_1, p_2, \dots, p_n . Let $Q = p_1 p_2 \dots p_n + 1$
- By Fundamental Theorem of Arithmetic: Q is prime or else it can be written as the product of two or more primes

Mersenne primes

- Primes with the special form 2^p-1 where p is also a prime, called **Mersenne prime**.
- $2^2-1=3$, $2^3-1=7$, $2^5-1=31$ are Mersenne primes while $2^{11}-1=2047$ is not a Mersenne prime ($2047=23 \cdot 89$)
- The largest Mersenne prime known (as of early 2011) is $2^{43,112,609}-1$, a number with over 13 million digits

Theorem 4

THE PRIME NUMBER THEOREM The ratio of the number of primes not exceeding x and $x/\ln x$ approaches 1 as x grows without bound. (Here $\ln x$ is the natural logarithm of x .)

- This theorem was proved in 1896 and proof is complicated.
- Can use this theorem to estimate the odds that a randomly chosen number is prime
- The odds that a randomly selected positive integer less than n is prime are approximately $(n/\ln n)/n=1/\ln n$
- The odds that an integer less than 10^{1000} is prime are approximately $1/\ln 10^{1000}$, approximately $1/2300$

Open Problems about Primes

- **Goldbach's conjecture:** every even integer n , $n > 2$, is the sum of two primes
 $4 = 2 + 2$, $6 = 3 + 3$, $8 = 5 + 3$, $10 = 7 + 3$, $12 = 7 + 5$, ...
- As of 2011, the conjecture has been checked for all positive even integers up to $1.6 \cdot 10^{18}$
- **Twin prime conjecture:** Twin primes are primes that differ by 2. There are infinitely many twin primes

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- **GCD and LCM**
- Euclidean Algorithm

Greatest common divisor

- Let a and b be integers, not both zero. The largest integer d such that $d \mid a$ and $d \mid b$ is called the **greatest common divisor** (GCD) of a and b , often denoted as $\gcd(a,b)$
- The integers a and b are **relative prime** if their GCD is 1
 $\gcd(10, 17)=1$, $\gcd(10, 21)=1$, $\gcd(10,24)=2$
- The integers a_1, a_2, \dots, a_n are **pairwise relatively prime** if $\gcd(a_i, a_j)=1$ whenever $1 \leq i < j \leq n$

Prime factorization and GCD

- Finding GCD

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}, b = p_1^{b_1} p_2^{b_2} \cdots p_n^{b_n}$$

$$\gcd(a, b) = p_1^{\min(a_1, b_1)} p_2^{\min(a_2, b_2)} \cdots p_n^{\min(a_n, b_n)}$$

$$120 = 2^3 \cdot 3 \cdot 5, \quad 500 = 2^2 \cdot 5^3$$

$$\gcd(120, 500) = 2^2 \cdot 3^0 \cdot 5^1 = 20$$

- **Least common multiples** of the positive integers a and b is the smallest positive integer that is divisible by both a and b , denoted as $\text{lcm}(a, b)$

Least common multiple

- Finding LCM

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}, b = p_1^{b_1} p_2^{b_2} \cdots p_n^{b_n}$$

$$\text{lcm}(a, b) = p_1^{\max(a_1, b_1)} p_2^{\max(a_2, b_2)} \cdots p_n^{\max(a_n, b_n)}$$

$$120 = 2^3 \cdot 3 \cdot 5, 500 = 2^2 \cdot 5^3$$

$$\text{lcm}(120, 500) = 2^3 \cdot 3^1 \cdot 5^3 = 8 \cdot 3 \cdot 125 = 3000$$

- Let a and b be positive integers, then
 $ab = \text{gcd}(a, b) \cdot \text{lcm}(a, b)$

Outline

- Prime and Composite
- Prime Factorizations
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- GCD and LCM
- **Euclidean Algorithm**

Euclidean algorithm

- Need more efficient prime factorization algorithm
- Example: Find $\gcd(91,287)$
- $287=91 \cdot 3 +14$
- Any divisor of 287 and 91 must be a divisor of $287- 91 \cdot 3 =14$
- Any divisor of 91 and 14 must also be a divisor of $287=91 \cdot 3$
- Hence, the $\gcd(91,287)=\gcd(91,14)$

Euclidean algorithm

- Need more efficient prime factorization algorithm
- Example: Find $\gcd(91, 287)$
- $\gcd(91, 287) = \gcd(91, 14)$
- Next, $91 = 14 \cdot 6 + 7$
- Any divisor of 91 and 14 also divides $91 - 14 \cdot 6 = 7$ and any divisor of 14 and 7 divides 91, i.e.,
 $\gcd(91, 14) = \gcd(14, 7)$
- $14 = 7 \cdot 2$, $\gcd(14, 7) = 7$,
- Thus $\gcd(287, 91) = \gcd(91, 14) = \gcd(14, 7) = 7$

Euclidean algorithm

Let $a = bq + r$, where a, b, q , and r are integers. Then $\gcd(a, b) = \gcd(b, r)$.

- Proof: Suppose that d divides both a and b . Then it follows that d also divides $a - bq = r$. Hence, any common divisor of a and b is also a common divisor of b and r .
- Likewise, suppose that d divides both b and r . Then d also divides $bq + r = a$. Hence, any common divisor of b and r is also a common divisor of a and b .
- Consequently, $\gcd(a, b) = \gcd(b, r)$

Euclidean algorithm

- Suppose a and b are positive integers, $a \geq b$. Let $r_0 = a$ and $r_1 = b$, we successively apply the division algorithm

$$r_0 = r_1 q_1 + r_2, 0 \leq r_2 < r_1$$

$$r_1 = r_2 q_2 + r_3, 0 \leq r_3 < r_2$$

...

$$r_{n-2} = r_{n-1} q_{n-1} + r_n, 0 \leq r_n < r_{n-1}$$

$$r_{n-1} = r_n q_n$$

$$\begin{aligned} \gcd(a, b) &= \gcd(r_0, r_1) = \gcd(r_1, r_2) = \cdots = \gcd(r_{n-2}, r_{n-1}) \\ &= \gcd(r_{n-1}, r_n) = \gcd(r_n, 0) = r_n \end{aligned}$$

- Hence, the gcd is the last nonzero remainder in the sequence of divisions

Example

- Find the GCD of 414 and 662

$$662 = 414 \cdot 1 + 248$$

$$414 = 248 \cdot 1 + 166$$

$$248 = 166 \cdot 1 + 82$$

$$166 = 82 \cdot 2 + 2$$

$$82 = 2 \cdot 41$$

$$\gcd(414, 662) = 2 \text{ (the last nonzero remainder)}$$

$$a = bq + r$$

$$\gcd(a, b) = \gcd(b, r)$$

The Euclidean algorithm

```
procedure gcd(a, b: positive integers)
x := a
y := b
while y ≠ 0
    r := x mod y
    x := y
    y := r
return x {gcd(a, b) is x}
```

- The time complexity is $O(\log b)$ (where $a \geq b$)

Next class

- Topic: Cryptograph
- Pre-class reading: Chap 5.6

